

Online Selection with Convex Costs

Xiaoqi Tan¹, Siyuan Yu¹, Raouf Boutaba²

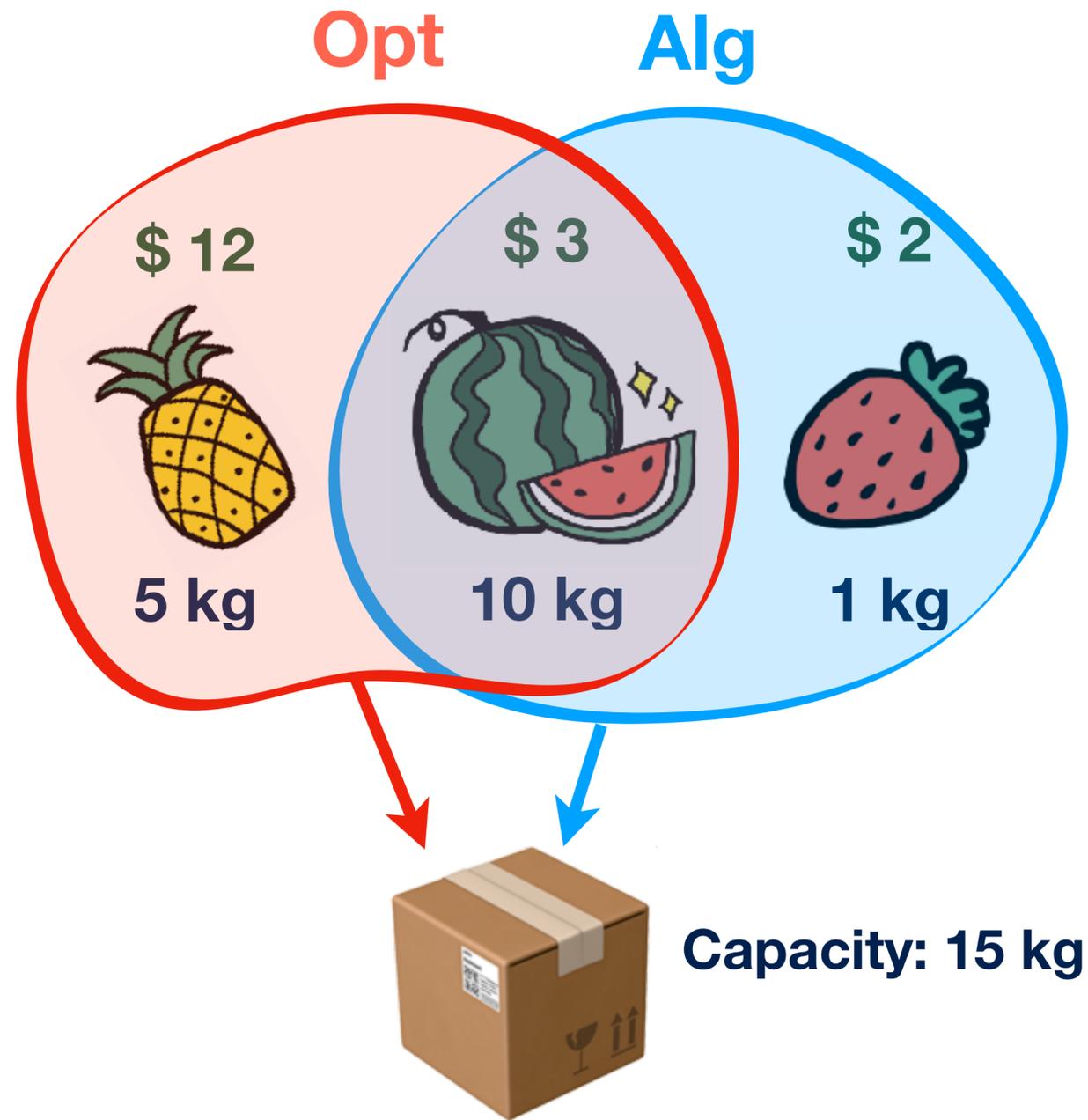
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MAMA Workshop

June 6, 2022

Example 1: Online Knapsack



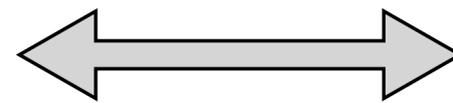
$$\frac{\text{Profit (Opt)}}{\text{Profit (Alg)}} = \frac{3 + 12}{2 + 3} = 3$$

Achieve **1/3** of the **best-possible profit** (in hindsight)

Competitive Ratio (CR)

- A **deterministic online algorithm** is α -competitive if

$$\text{Alg}(\sigma) \geq \frac{1}{\alpha} \cdot \text{Opt}(\sigma), \quad \forall \sigma$$



$$\alpha = \max_{\sigma} \frac{\text{Opt}(\sigma)}{\text{Alg}(\sigma)}$$

- A **randomized online algorithm** is α -competitive if

$$\mathbb{E}[\text{Alg}(\sigma)] \geq \frac{1}{\alpha} \cdot \text{Opt}(\sigma), \quad \forall \sigma$$

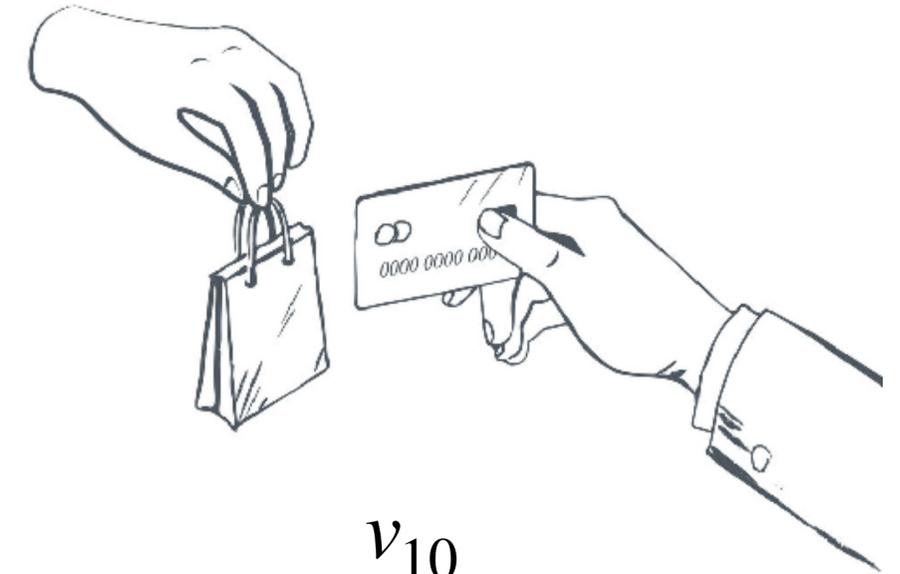


$$\alpha = \max_{\sigma} \frac{\text{Opt}(\sigma)}{\mathbb{E}[\text{Alg}(\sigma)]}$$

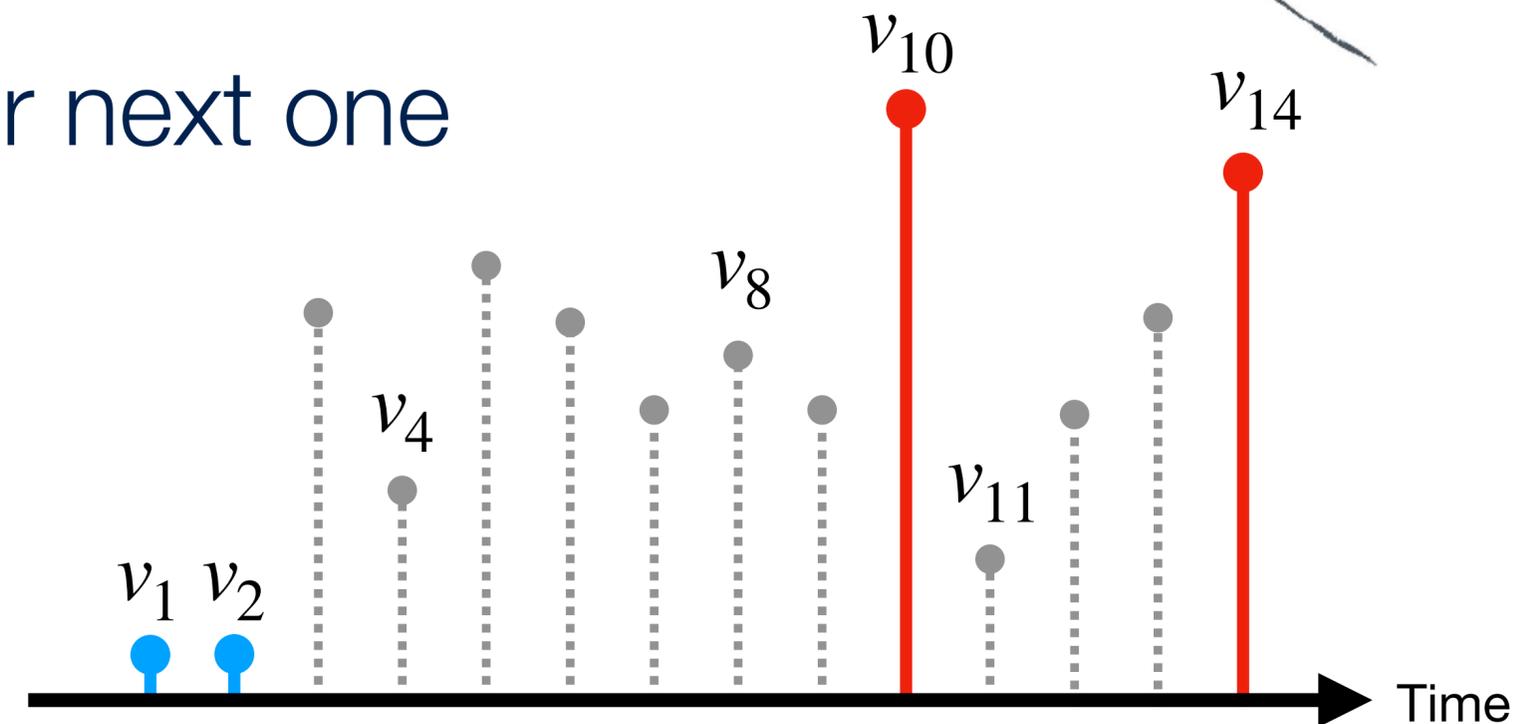
Worst-case competitive analysis

Example 2: k -max Search

- **Selling** k items for a sequence of buyers
- Each buyer $t = 1, 2, \dots, T$ makes an **offer** v_t
- **Accept it and sell 1 item**, or wait for next one
- Once sold, cannot be reclaimed



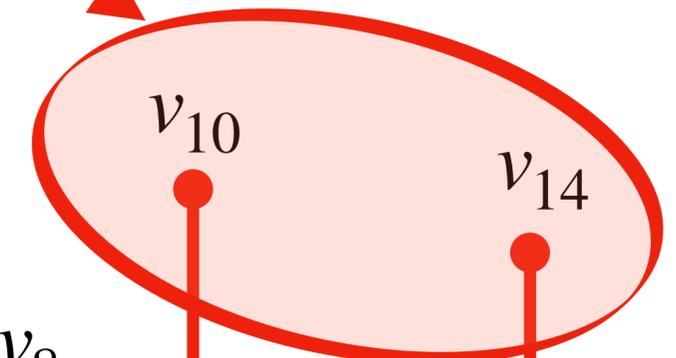
$$\frac{\text{Profit (Opt)}}{\text{Profit (Alg)}} = \frac{v_{10} + v_{14}}{v_1 + v_2} > 1$$



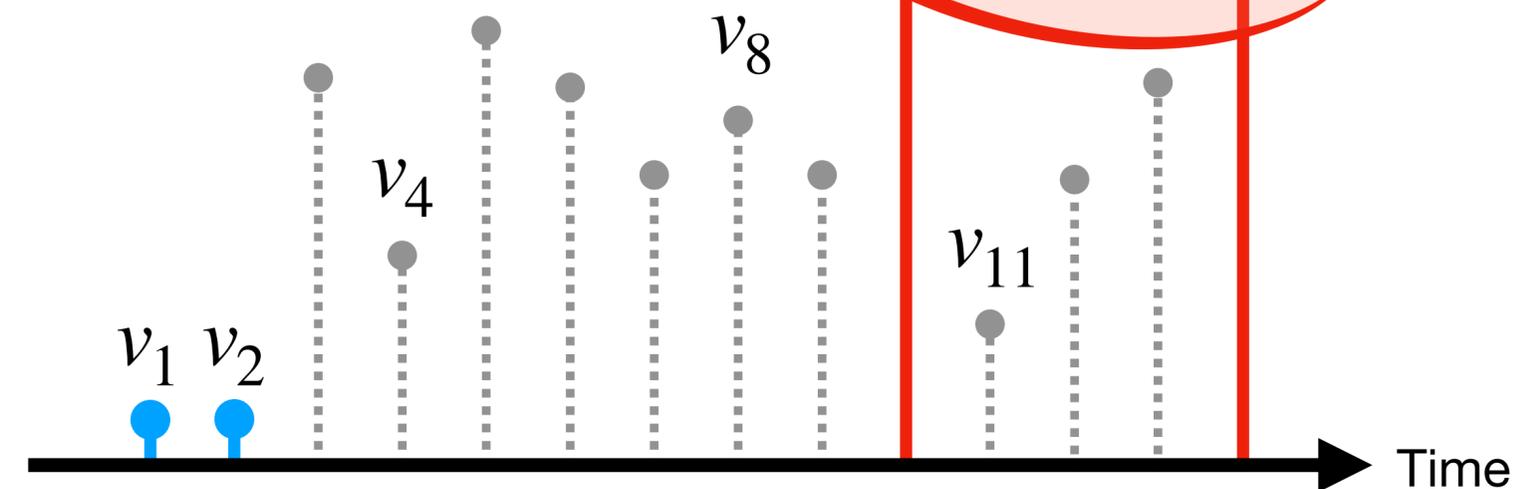
e.g., $k = 2, T = 14$

$$\begin{array}{l}
 \text{maximize} \\
 x_t \quad \sum_{t=1}^T v_t x_t \\
 \text{subject to} \quad \sum_{t=1}^T x_t \leq k, \\
 x_t \in \{0,1\}, \forall t
 \end{array}$$

Assume all future knowledge



$$\frac{\text{Profit (Opt)}}{\text{Profit (Alg)}} = \frac{v_{10} + v_{14}}{v_1 + v_2} > 1$$



e.g., $k = 2, T = 14$

Example 3: One-Way Trading

Algorithmica 2001

Optimal Search and One-Way Trading
Online Algorithms

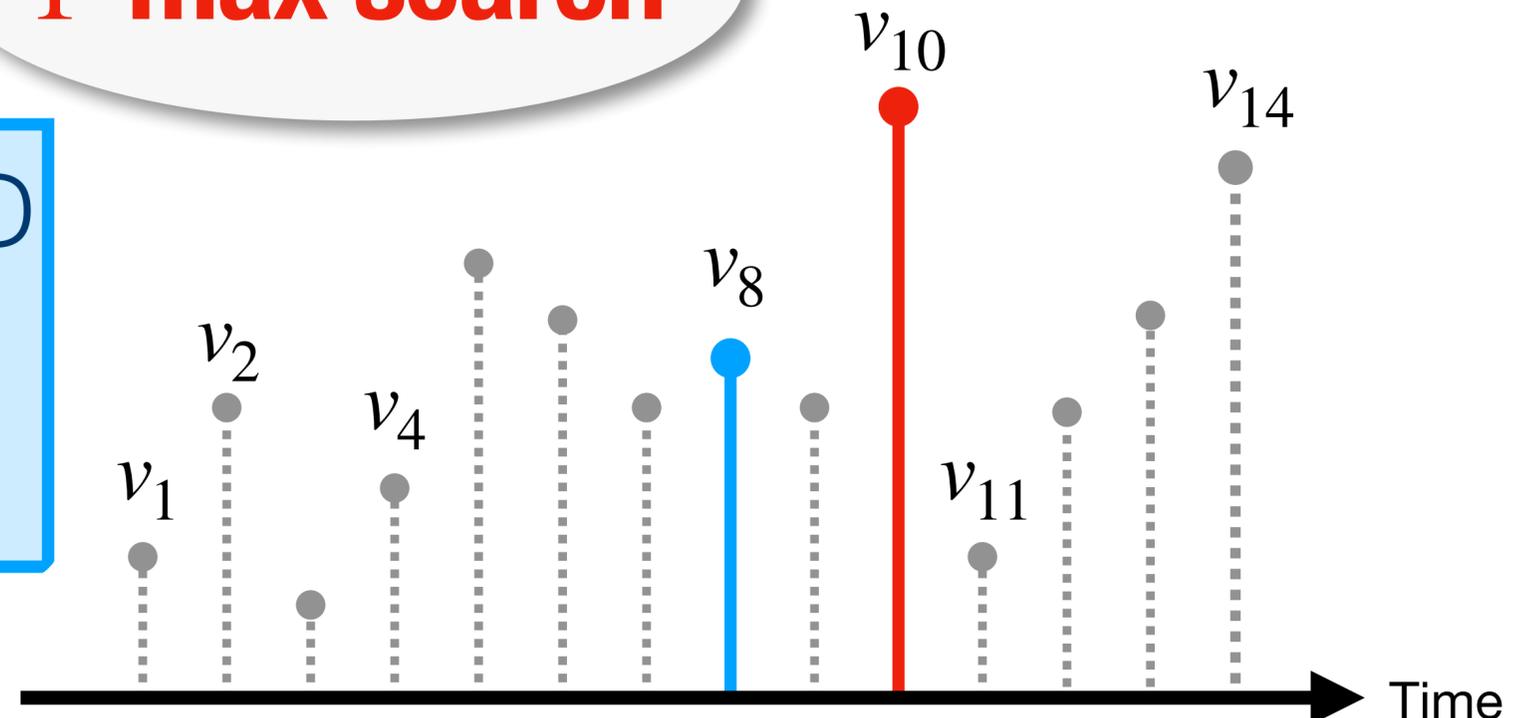
R. El-Yaniv,¹ A. Fiat,² R. M. Karp,³ and G. Turpin⁴

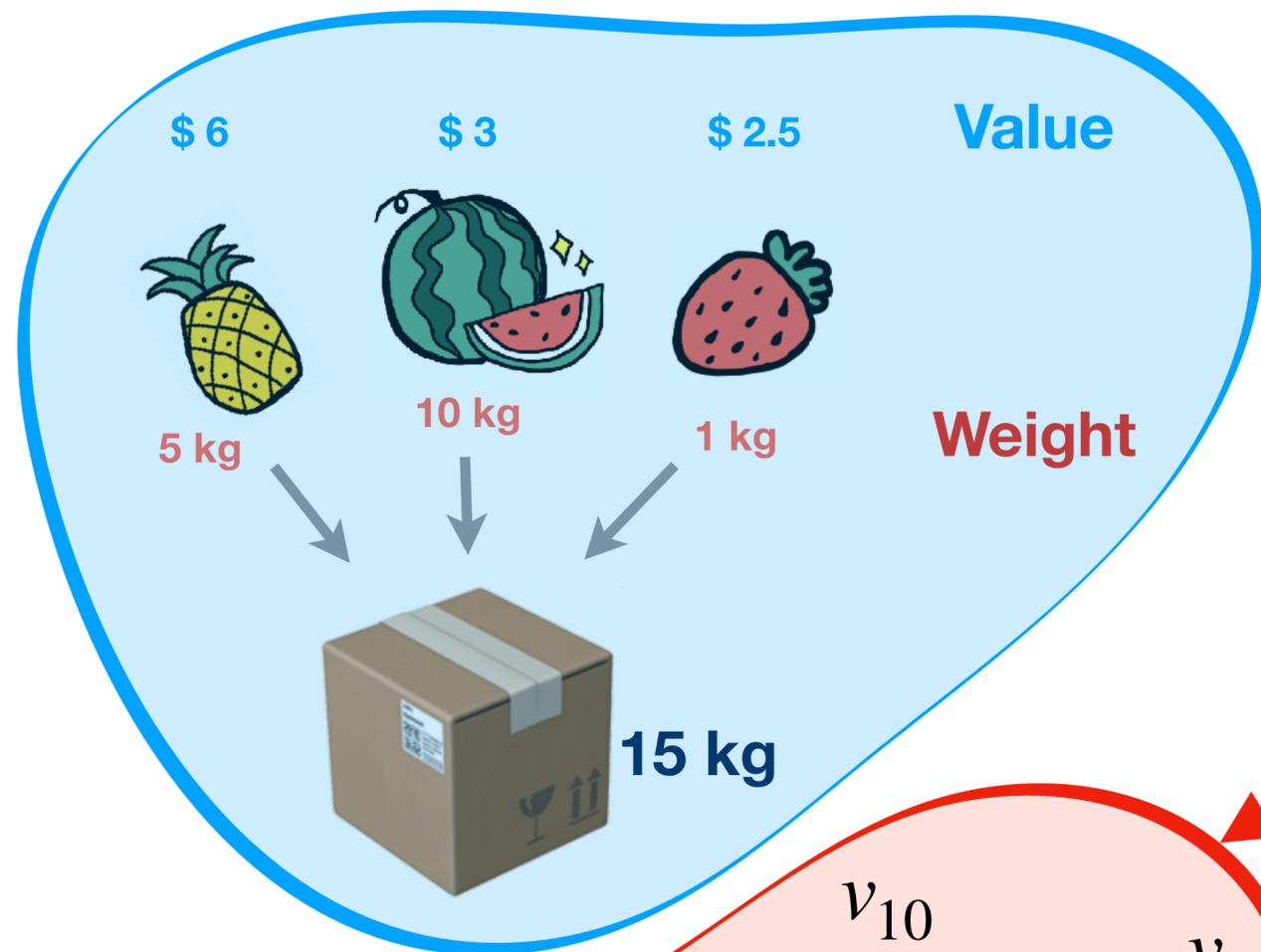
- To change 1 CAD for USD.
- Rates v_t are revealed sequentially for $t = 1, 2, \dots, T$

- At step t : Trade x_t CAD for USD
- **Maximize** USD in the end.

Many variants...

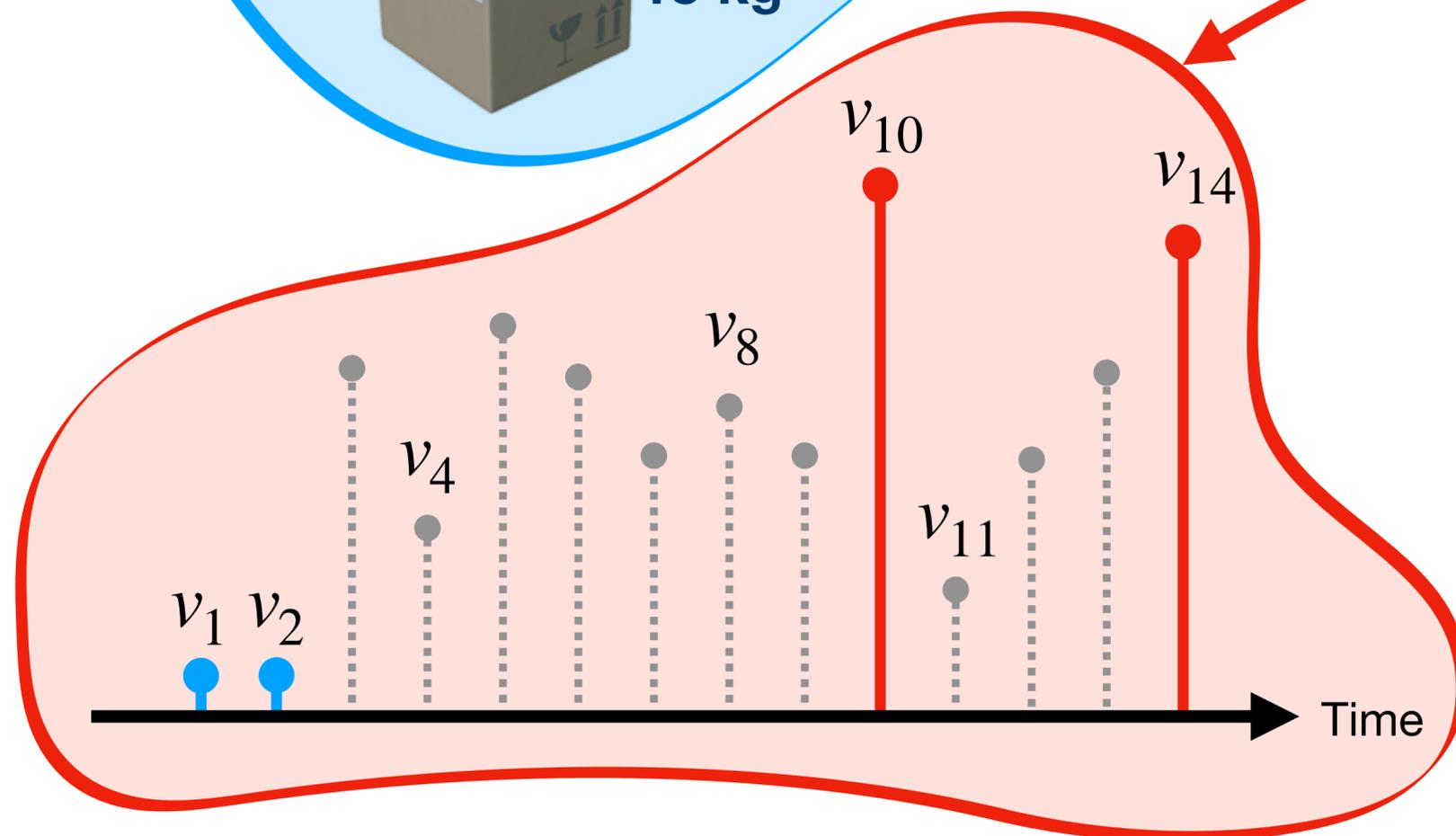
1-max search





May have multi-dimensional attributes (e.g., value and weight)

A 1-d scale value



- A set of T items (one at a time)

- Select a **subset** $S \subset T$ of items with possible **constraints**

- Maximize **objective** $v(S)$

NeurIPS 2021

Pareto-Optimal Learning-Augmented Algorithms for Online Conversion Problems

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Algorithmica 2001

Optimal Search and One-Way Trading Online Algorithms

R. El-Yaniv,¹ A. Fiat,² R. M. Karp,³ and G. Turpin⁴

Online selection problems

- A set of T items (one at a time)
- Select a **subset** $S \subset T$ of items with possible **constraints**
- Maximize **objective** $v(S)$

Online Selection Problems against Constrained Adversary

Zhihao Jiang¹ Pinyan Lu² Zhihao Gavin Tang² Yuhao Zhang³

ICML 2021

An Efficient Framework for Balancing Submodularity and Cost *

Sofia Maria Nikolakaki

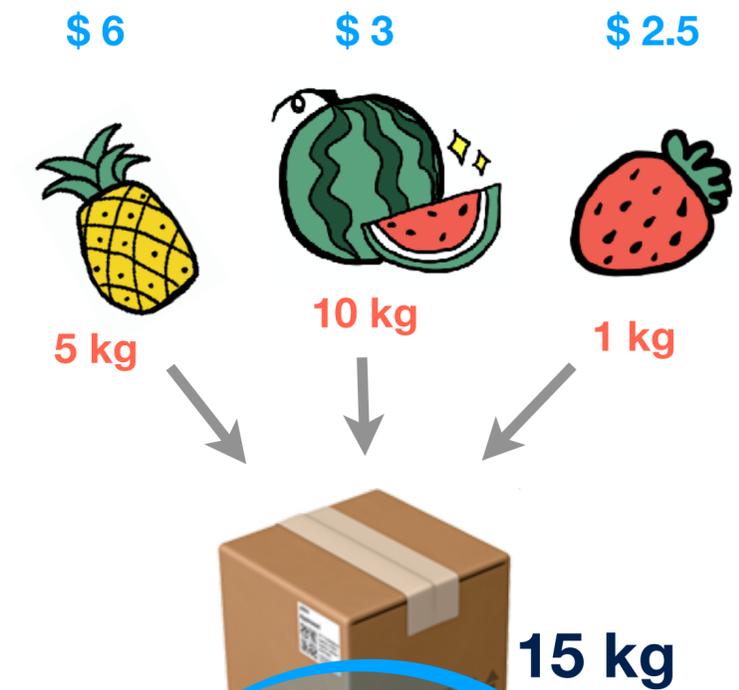
Alina Ene

Evimaria Terzi

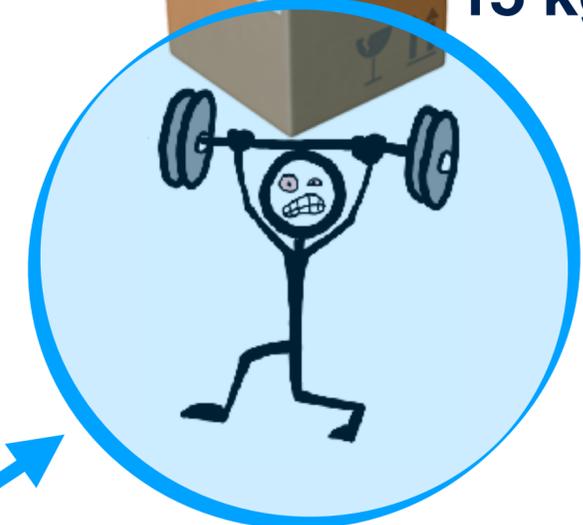
SIGKDD 2021

A New Variant: OSCC

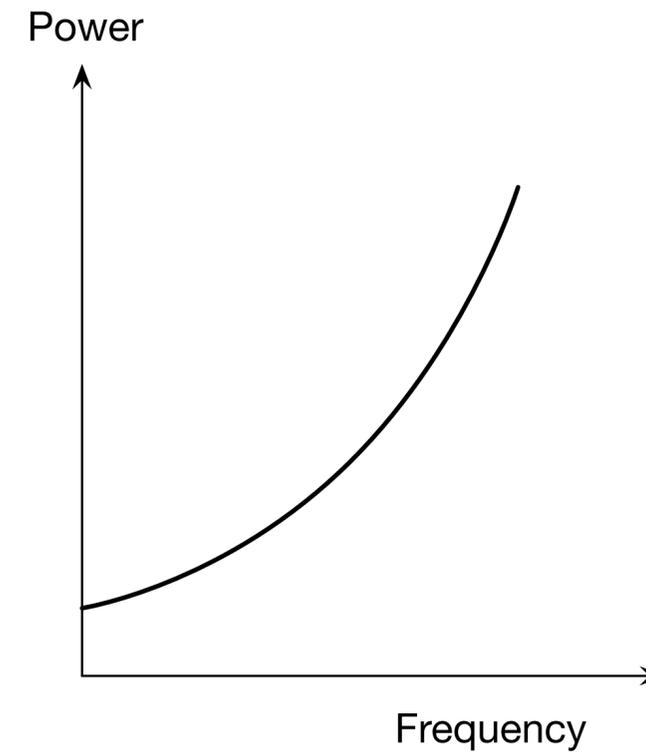
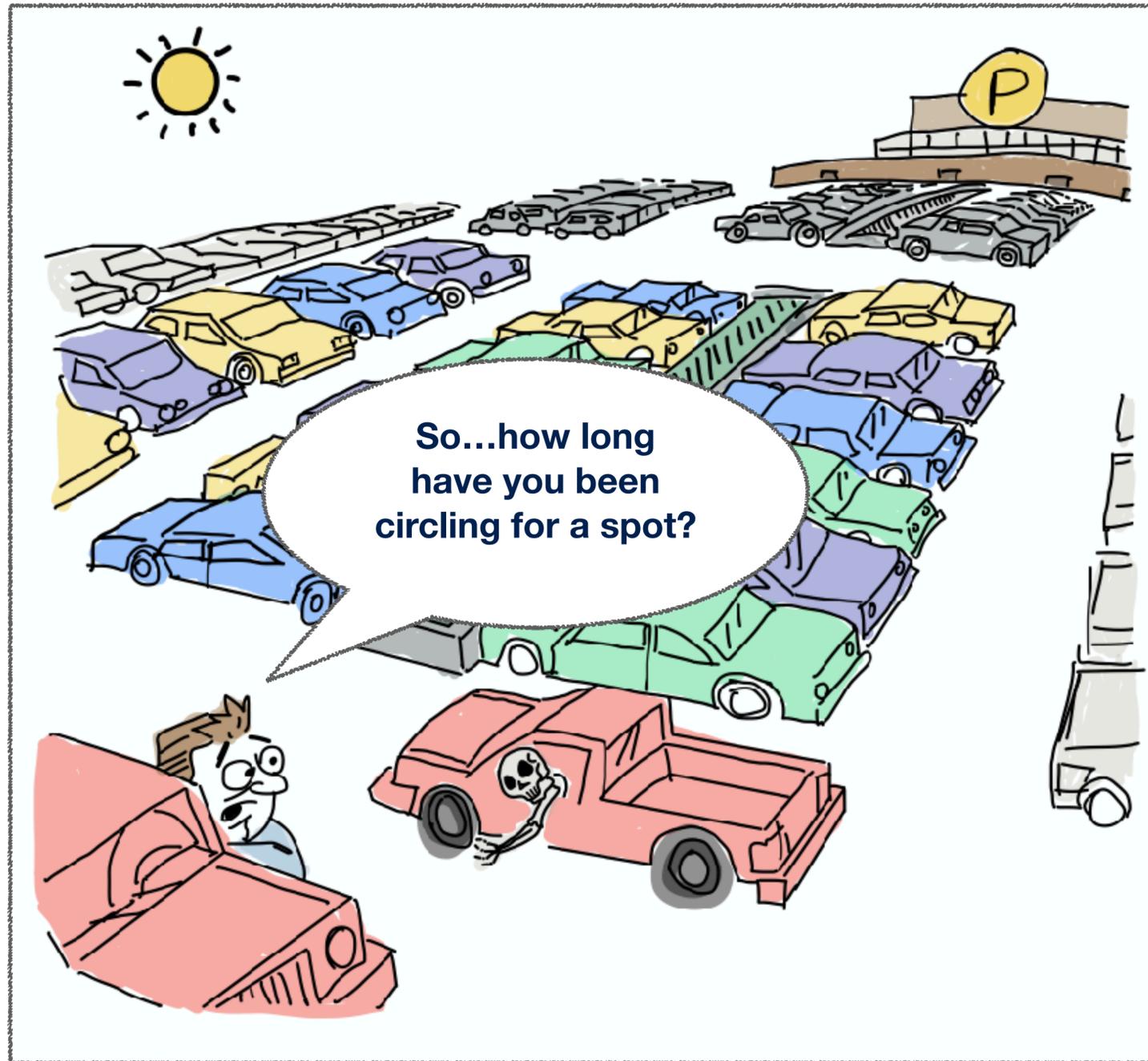
- A set of T items (one at a time)
- Select a subset $S \subset T$ of items (with possible constraints)
- Maximize objective $v(S) = f(S)$



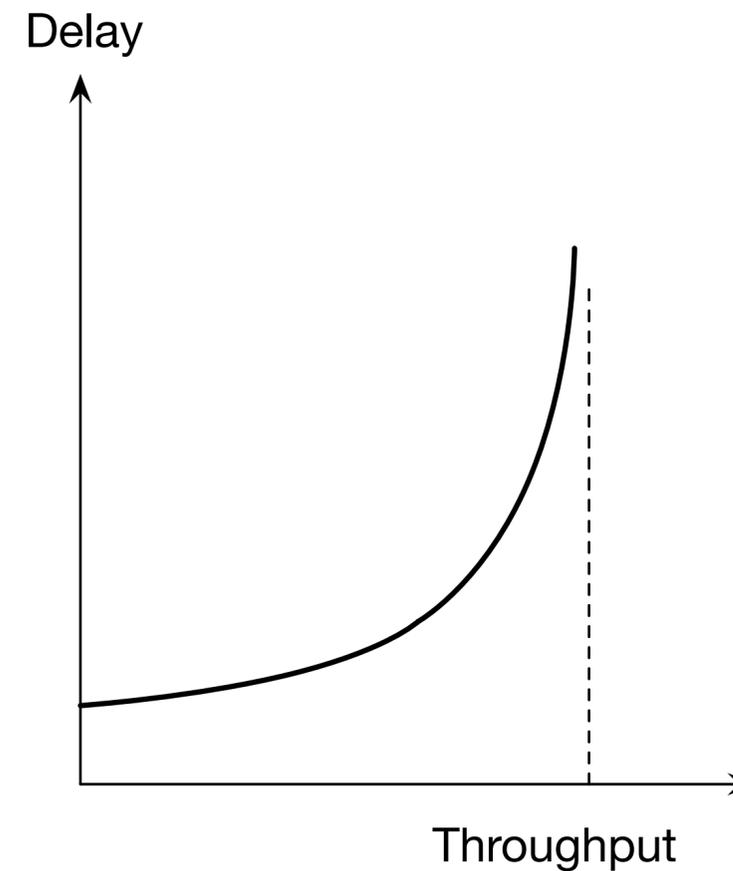
OSCC: Online selection with convex costs



Convex & increasing costs



**Cloud
resource
allocation**



**Network
resource
allocation**

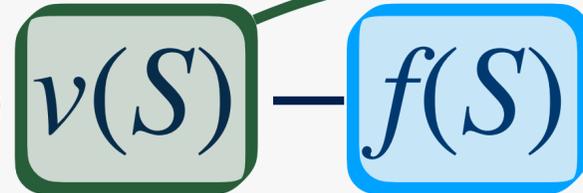
OSCC: Value-Cost Tradeoff

- A set of T items (one at a time)
- Select a subset $S \subset T$ of items (with possible constraints)
- Maximize objective $v(S) - f(S)$



OSCC and More Variants

- A set of T items (one at a time)
- Select a subset $S \subset T$ of items (with possible constraints)
- Maximize objective $v(S) - f(S)$



Multi-dimensional?

Combinatorial?

**Linear?
Separable?**

**Convex?
Separable?**

Welfare Maximization with Production Costs:
A Primal Dual Approach

Zhiyi Huang*
The University of Hong Kong
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Anthony Kim†
Stanford University
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**Games and Economic
Behavior (2018)**

**Mechanism Design for Online Resource Allocation: A
Unified Approach**

XIAOQI TAN, University of Toronto, Canada
BO SUN, HKUST, Hong Kong, China
ALBERTO LEON-GARCIA, University of Toronto, Canada
YUAN WU, University of Macau, Macau, China
DANNY H.K. TSANG, HKUST, Hong Kong, China

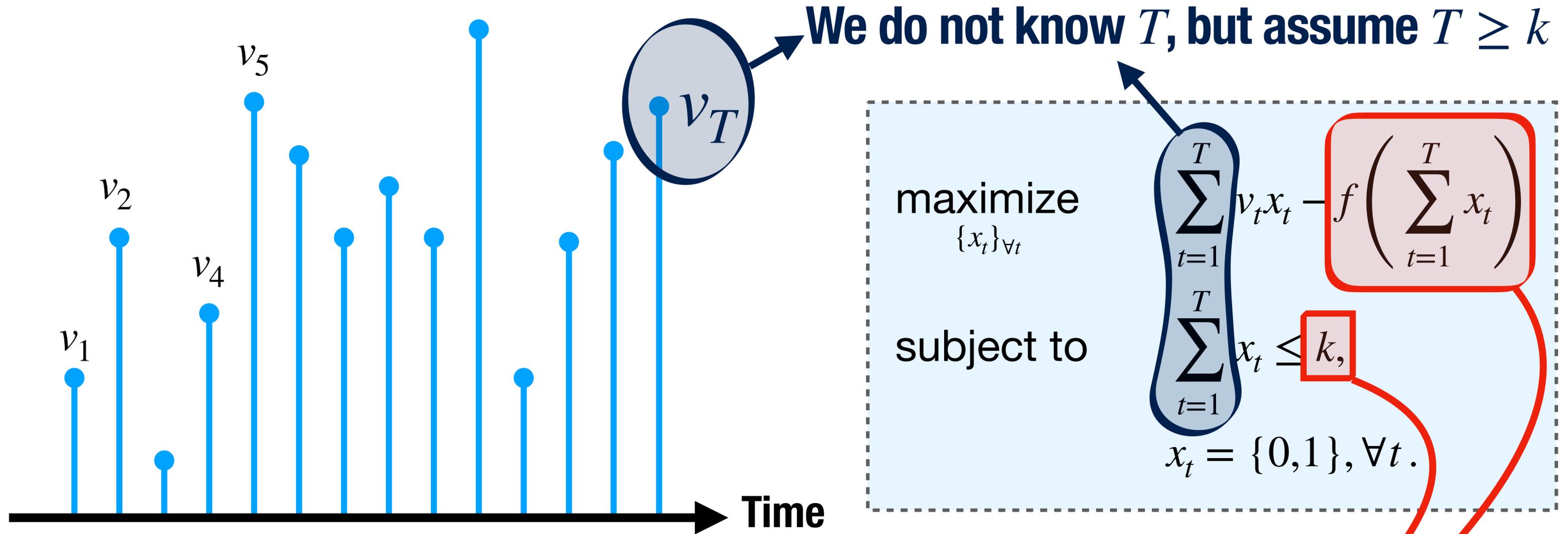
SIGMETRICS 2020

Online Combinatorial Auctions for Resource
Allocation with Supply Costs and Capacity Limits

Xiaoqi Tan, Alberto Leon-Garcia, Yuan Wu, and Danny H.K. Tsang

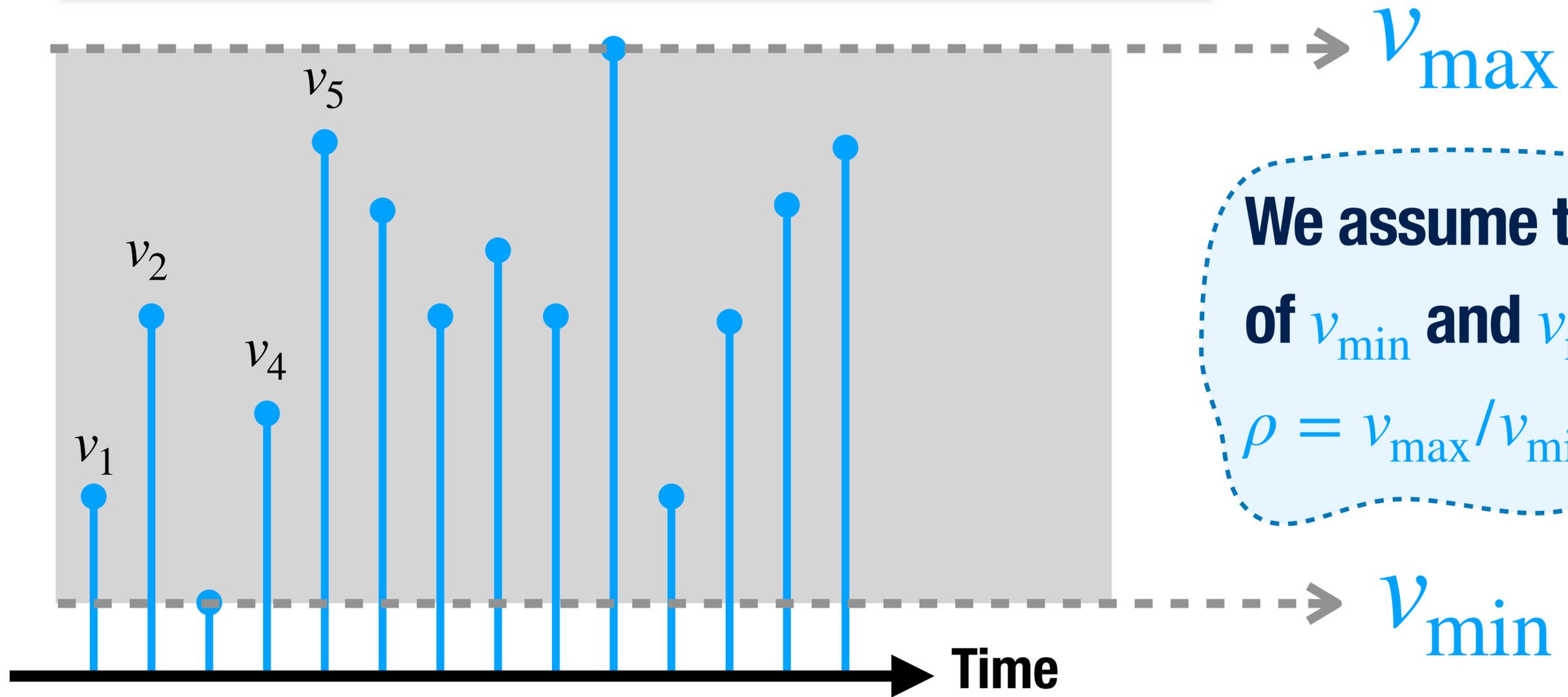
IEEE-JSAC 2020

OSCC: Basic Setting (For Today)



In this setting: OSCC = k -max Search with Convex Costs

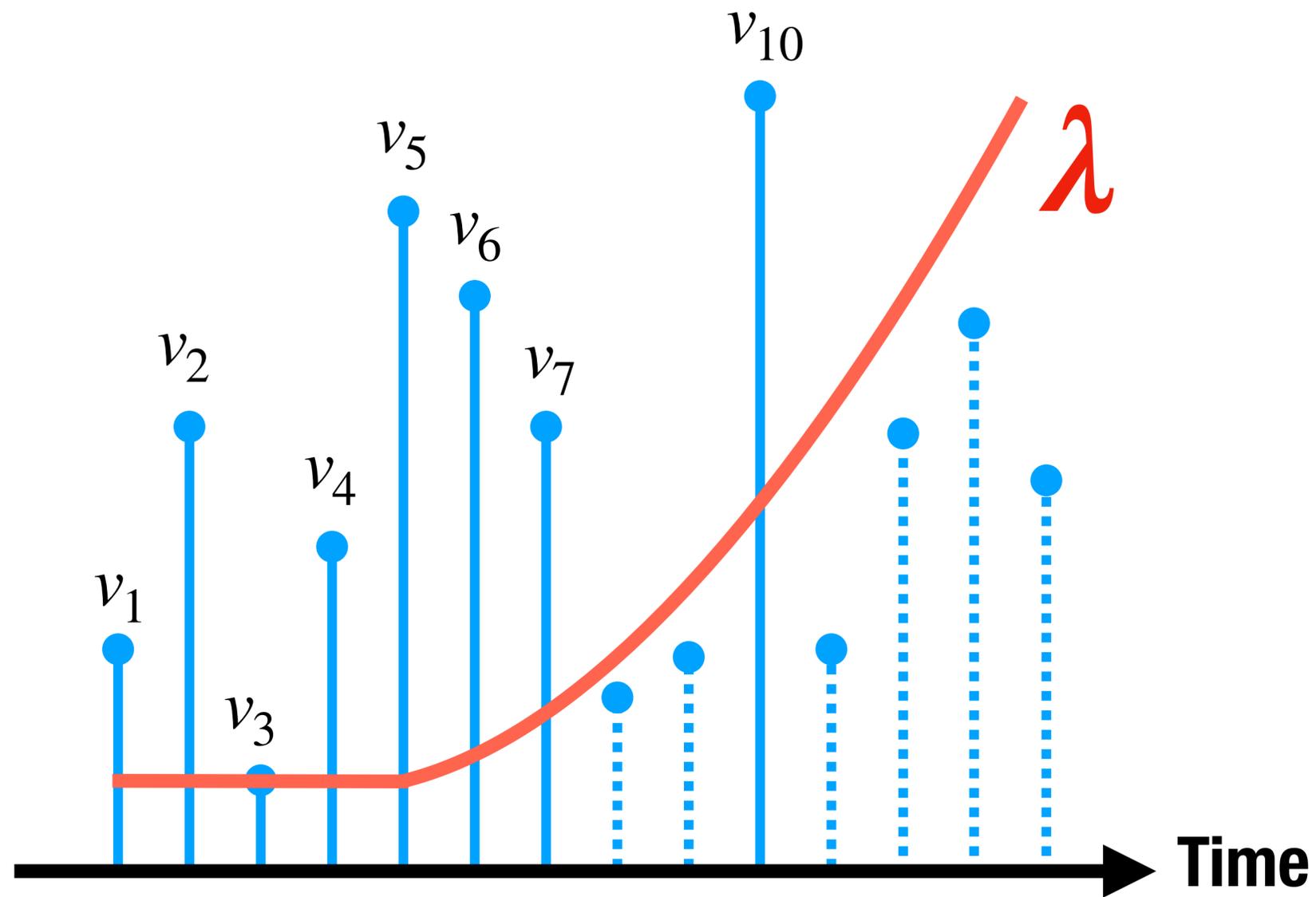
Larger $\rho \implies$ Higher levels of uncertainty



We assume the knowledge of v_{\min} and v_{\max} , and define

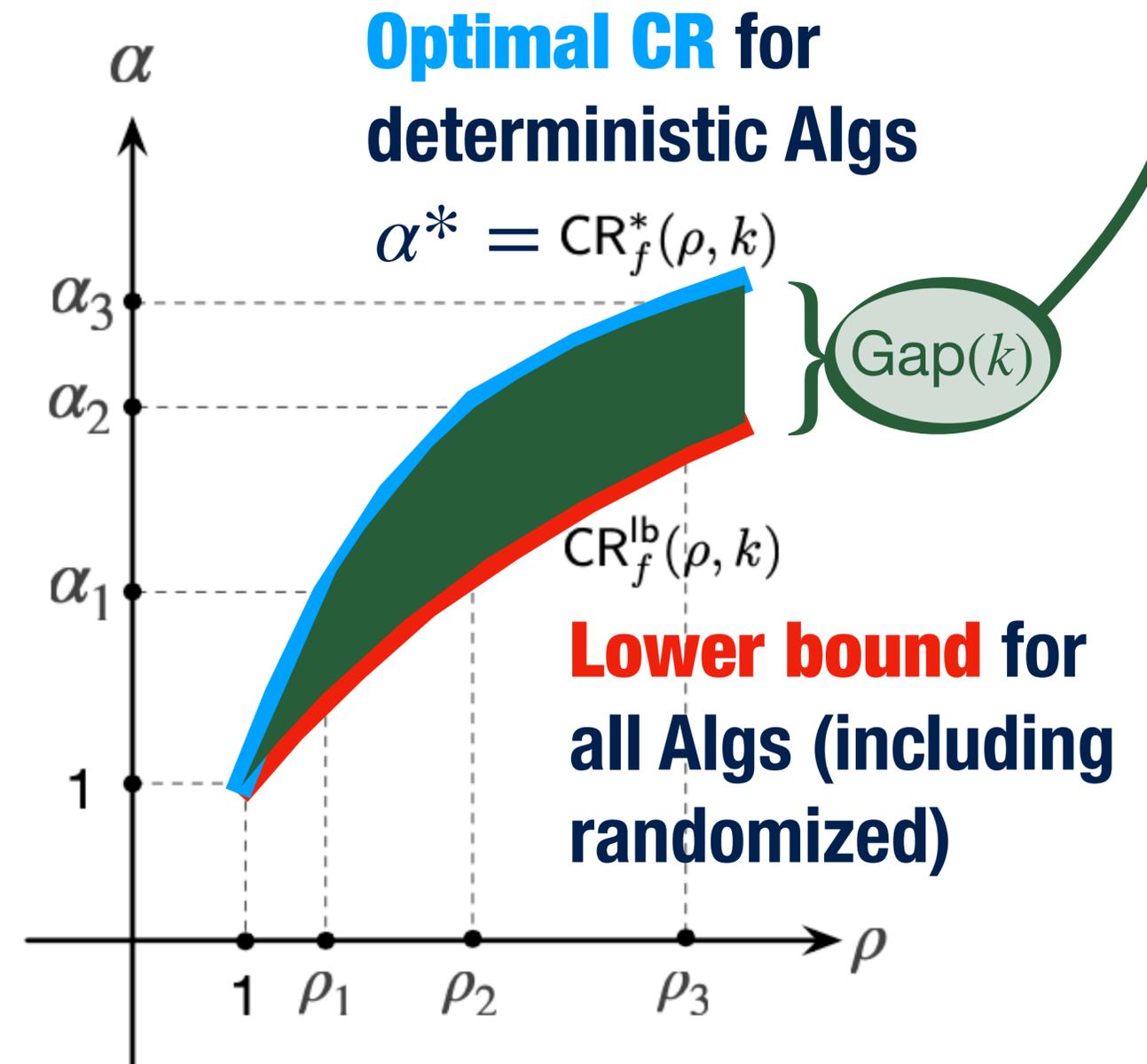
$$\rho = v_{\max} / v_{\min}$$

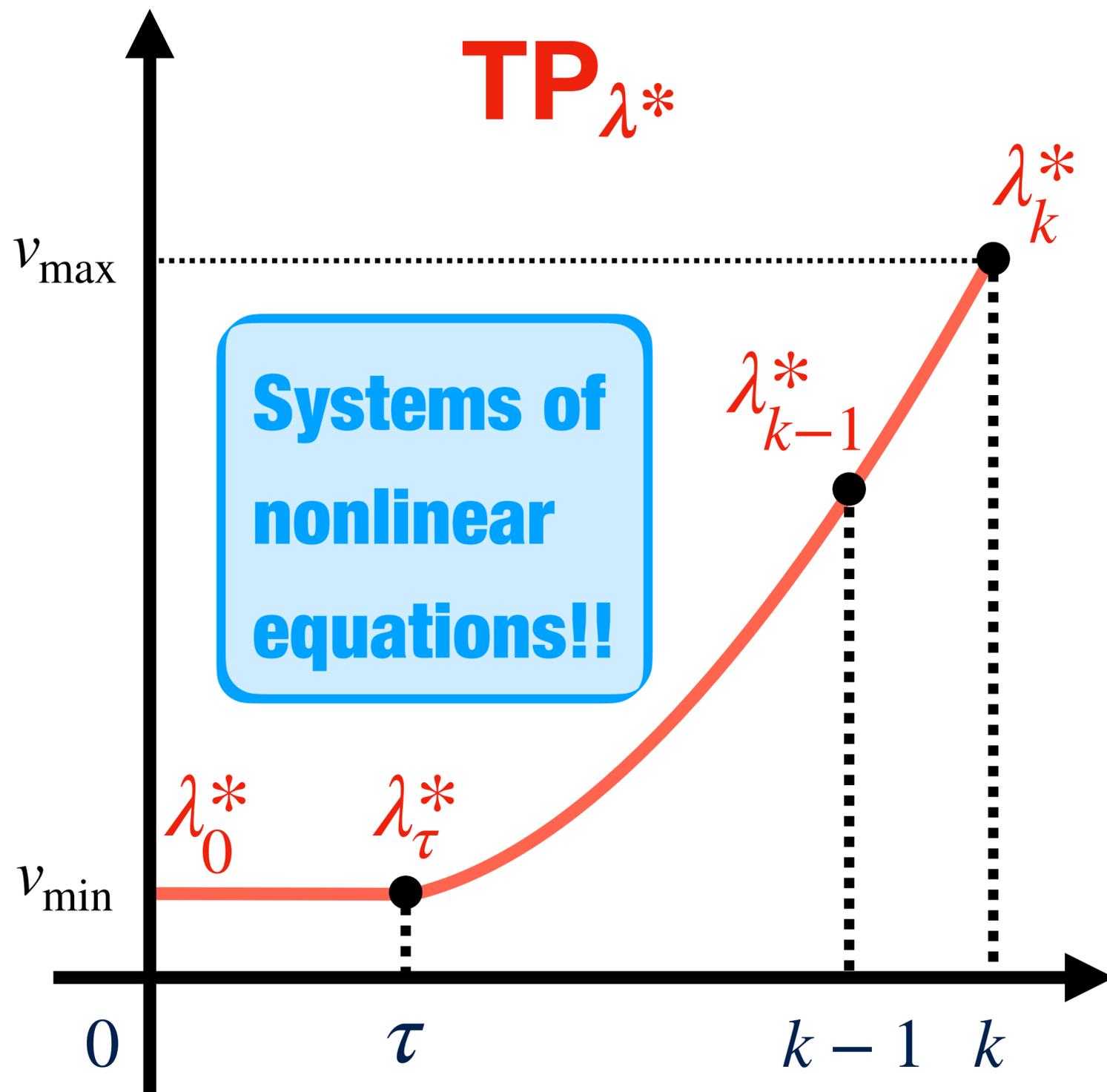
Our Algorithm: TP_λ



$$\rho = v_{\max} / v_{\min}$$

$$\lim_{k \rightarrow \infty} \text{Gap}(k) = 0$$





Theorem 1. TP_{λ^*} achieves the best-possible CR of all deterministic algorithms, denoted by α^* , if and only if $\lambda^* = \{\lambda_0^*, \lambda_1^*, \dots, \lambda_\tau^*, \dots, \lambda_k^*\}$ is an admission threshold such that

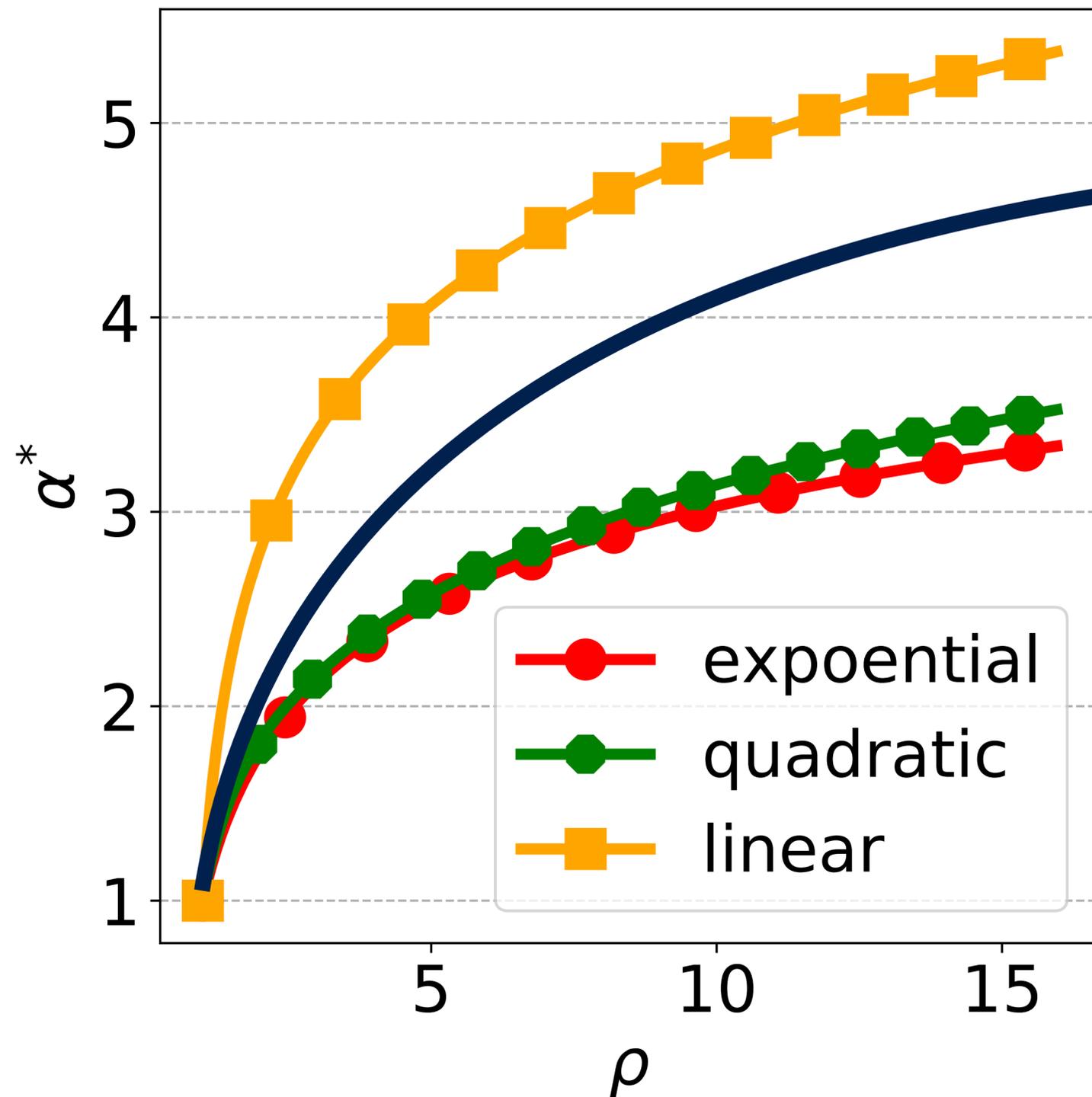
- $\lambda_0^* = \lambda_1^* = \dots = \lambda_\tau^* = v_{\min}$ and $\lambda_k^* = v_{\max}$, where τ is the minimum integer in $\{0, 1, \dots, k-1\}$ such that

$$v_{\min}(\tau + 1) - f(\tau + 1) \geq \frac{f^*(v_{\min})}{\alpha^*}. \quad (5)$$

- $\{\alpha^*, \lambda_{\tau+1}^*, \lambda_{\tau+2}^*, \dots, \lambda_{k-1}^*, \lambda_k^*\}$ is the unique set of $k - \tau + 1$ positive real numbers that satisfy the system of equations:

$$\alpha^* = \frac{f^*(\lambda_{\tau+1}^*)}{v_{\min}(\tau + 1) - f(\tau + 1)} = \frac{f^*(\lambda_{\tau+2}^*) - f^*(\lambda_{\tau+1}^*)}{\lambda_{\tau+1}^* - c_{\tau+2}} = \dots = \frac{f^*(\lambda_k^*) - f^*(\lambda_{k-1}^*)}{\lambda_{k-1}^* - c_k}. \quad (6)$$

- Existence** of many competitive TPs
- Uniqueness** of optimal (deterministic) TP
- Optimality** among all deterministic algs



$$1 + \ln \rho$$

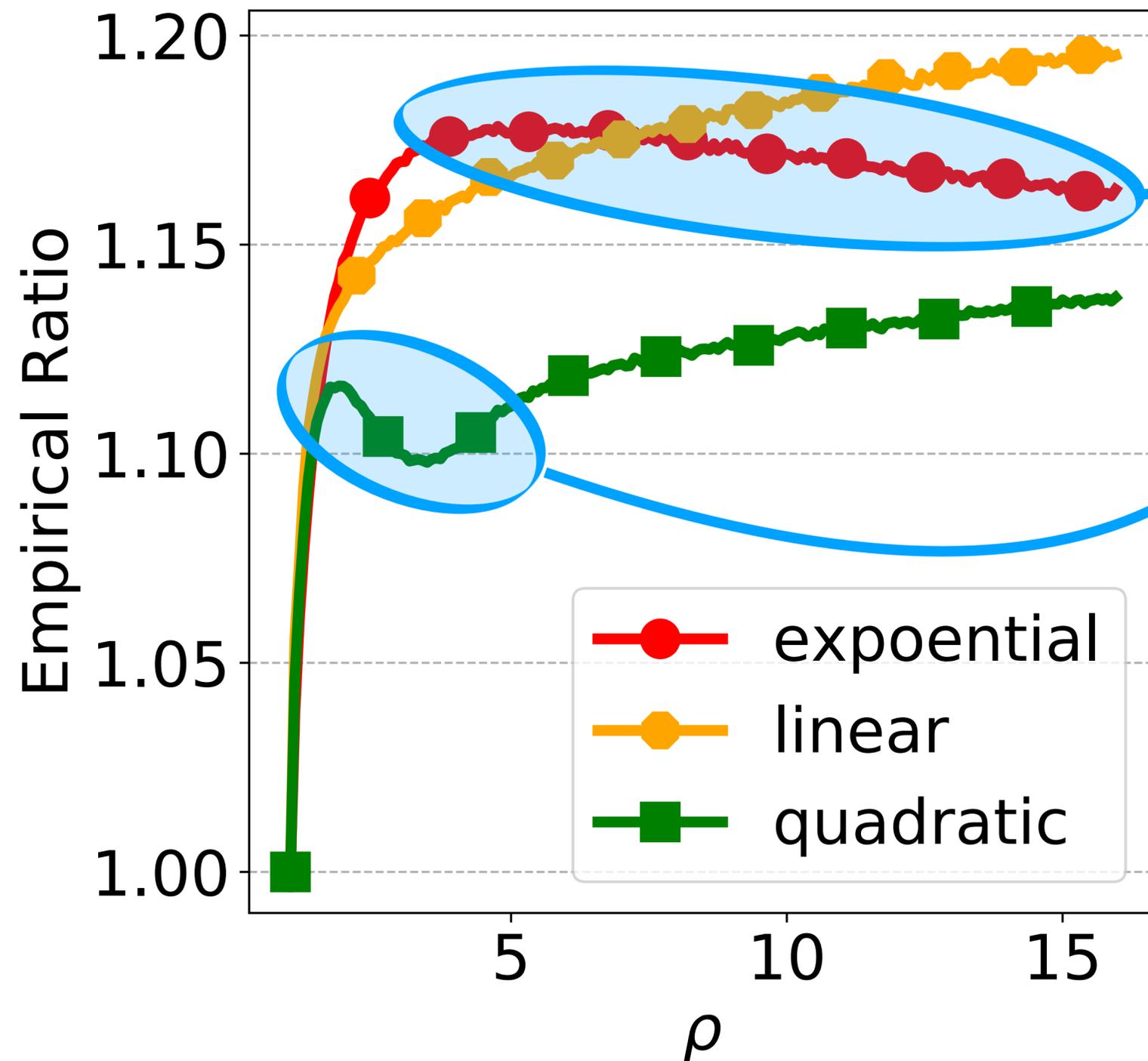
Strongly convex

Higher level of uncertainty \implies
worse performance (larger α^*)

Faster growth rate of cost \implies
better performance (smaller α^*)

Worst-case performance

$$\text{Empirical Ratio} \triangleq \frac{1}{N} \sum_{n=1}^N \frac{\text{OPT}(\mathcal{I}_n)}{\text{TP}_{\lambda^*}(\mathcal{I}_n)}$$



Higher level of uncertainty does not always lead to **worse (empirical) average performance!!**

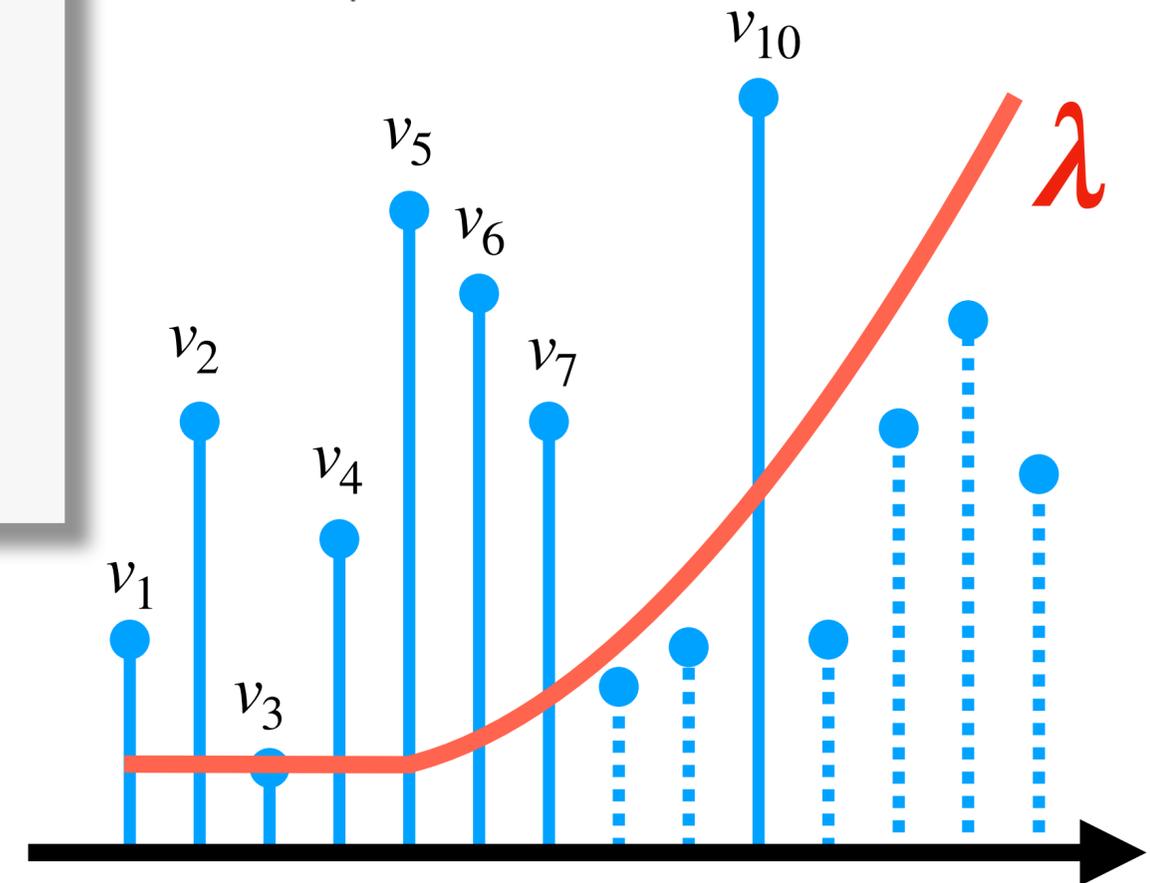
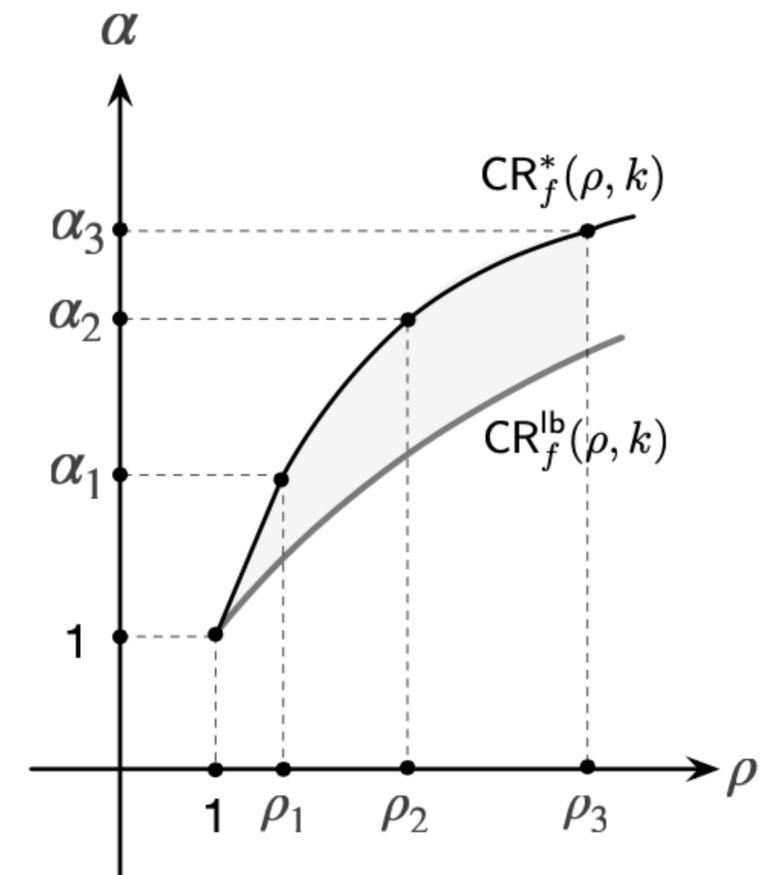
Need for more research:

- Flaws of worst-case analysis -> **beyond worst-case??**
- Implication of different costs on **real-world systems design??**

Concluding Remarks

- A new variant of online optimization problem: **online selection with convex costs**
- **Simple & intuitive** algorithms, **strong** performance guarantees
- **Interesting applications & open questions**

Thank You



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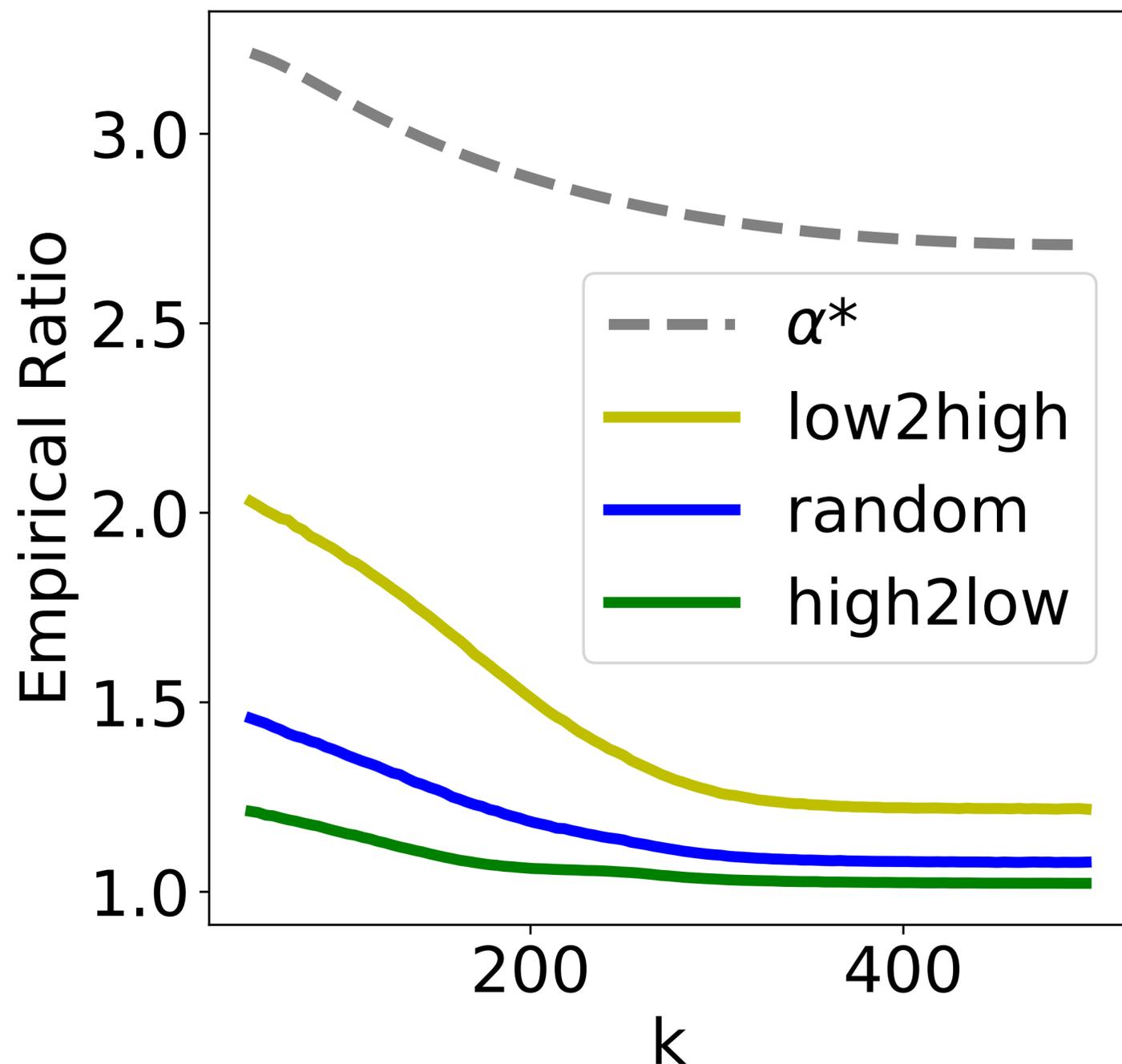
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$$\text{Empirical Ratio} \triangleq \frac{1}{N} \sum_{n=1}^N \frac{\text{OPT}(\mathcal{I}_n)}{\text{TP}_{\lambda^*}(\mathcal{I}_n)}$$

Different arrival instances



- **low2high**: 1st half (250 items) low values (randomly sampled from $[v_{\min}, (v_{\min} + v_{\max})/2]$) and 2nd half high
- **high2low**: 1st half (250 items) high values (randomly sampled from $[(v_{\min} + v_{\max})/2, v_{\max}]$) and 2nd half low
- **random**: all 500 items are randomly sampled from $[v_{\min}, v_{\max}]$